

781

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Note Title

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Predicate logic (or: First-order logic  
(FOL)  
or First-order predicate logic  
or Predicate Calculus)

$P_i^k$   
 $k$  is the arity of a predicate (symbol)  
arity is the number of arguments

Terms  
fx. 43 matrix:  $(Q(x) \vee P(f(x), z) \wedge Q(a)) \wedge$   
 $R(x, z, q(x))$

Shortcut:

$$I_a(x) = x^a$$

$$I_a(f) = f^a$$

$$I_a(p) = p^a$$

$$I_a(b) = b^a$$

$P(x, f(x))$

[www.possessh.org](http://www.possessh.org)

$$T_a(P(x, f(x))) = \text{student}$$

(coupon code)

$$= P^a(x^a, f^a(x^a)) =$$

$$= \langle 2, \text{succ}(2) \rangle = \langle 2, 3 \rangle$$

$$\therefore 2 < 3 = 0(f)$$

If  $F$  has the form  $f = \forall x \varphi$ , then

$$\alpha(F) = \begin{cases} 1, & \text{if for all } u \in U_a, \alpha_{[x/u]}(\varphi) = 1 \\ 0, & \text{otherwise} \end{cases}$$

If  $f$  has the form  $f = \exists x \varphi$ , then

$$\alpha(F) = \begin{cases} 1, & \text{if for some } u \in U_a, \alpha_{[x/u]}(\varphi) = 1 \\ 0, & \text{otherwise} \end{cases}$$

$\alpha \models F$  iff  $\alpha(F) = 1$

If  $\alpha(F) = 1$  for every suitable structure  $\alpha$ ,  
then  $F$  is valid, written  $\vdash F$

If there is some suitable structure  $\alpha$   
for which  $\alpha(F) = 0$ , then  $F$  is  
unsatisfiable.

If there is no model for  $F$  (i.e., no  
suitable structure  $\alpha$  s.t.  $\alpha(F) = 1$ ),

then  $f$  is unsatisfiable (or a contradiction).

Exercise 44

$$F: \forall x \exists y P(x, y, f(z))$$

A model for  $F$  is  $\alpha(V_\alpha, I_\alpha)$ , where

$$V_\alpha = \{x\}$$

$$f^\alpha = c \rightarrow c$$

$$\begin{aligned} z^\alpha &= c \\ P^\alpha &= \{(x, c, c)\} \end{aligned}$$

On interpretation  $\mathcal{B}$  s.t.  $\mathcal{B}(f) = 0$

$$U_{\mathcal{B}} = \{ \sim \}$$

$$f^{\mathcal{B}} = c \rightarrow c$$

$$\geq = c$$

$$P^{\mathcal{B}} = \{ \}$$