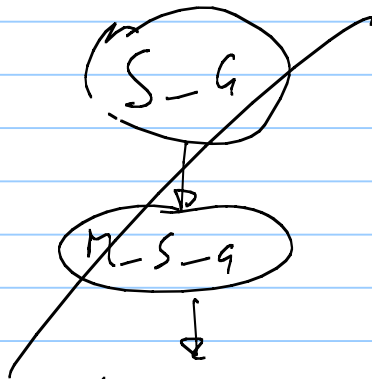




Section 4.3

A node is simplicial if its family is a clique (i.e., if it has a complete neighbor set).

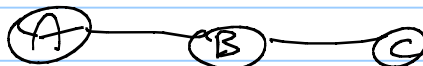
Proposition 4.4: Let G be a triangulated graph, x be a simplicial node. Let G' be G with x eliminated. Then G' is triangulated. (This is just Prop. 4.2 stated using new definitions.)

Theorem 4.1 A triangulated graph with at least two nodes has at least two simplicial nodes.

With two nodes, $(A) - (B)$, both A and B are simplicial. Trivial also for complete graphs.

stat An ^{incomplete} triangulated graph with at least three nodes has at least two non-adjacent simplicial nodes.

By induction.

Base case: 

Ind. step. Assume stat is true for graphs w/ fewer than n nodes.

We need to show that any incomplete triang. graph of n nodes ^(G or G') has two non-adj. simplicial nodes.

Let x be the first node in the ^{perfect} elimination sequence (which exists b/c G is triangulated). Let G' be G .

after eliminating x , G' is triangulated (b/c G

has a perfect elim ordering and we removed the first node in the ordering).

Two cases: G' is complete — choose any other node

G' is not complete. Then it meets the 2nd type, i.e., G' is a non-complete triangulated graph of $< n$ (exactly $n-1$) nodes. ✓

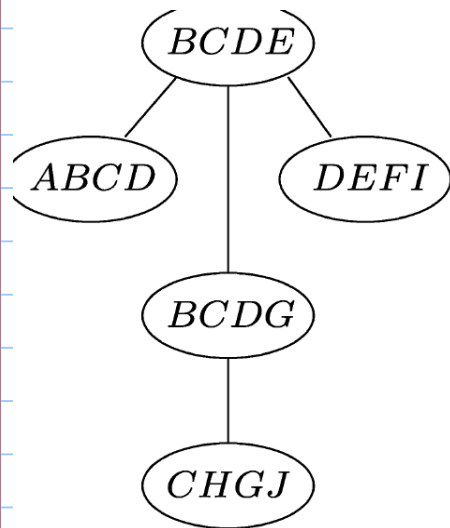
Corollary 4.1 In a triangulated graph, each variable A has a perfect elimination sequence ending in A .

Test for whether a graph G is triangulated is constructive: just try eliminating simplicial nodes — if you can obtain the empty graph, then G is triangulated; otherwise it is not.

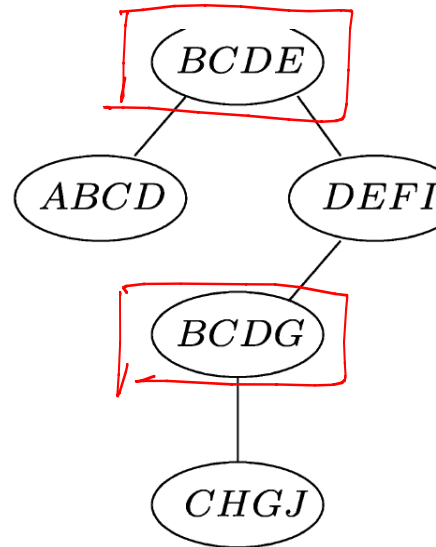
Alg. 4.1 To determine the set of cliques of a triang. graph do the following. (Read it!)

Join trees (often called junction trees, but in [507] junction trees are more complicated entities)

A join tree ^{for a triangulated graph G} is a tree $T = (V, E)$, where V is the set of cliques of G and, if there is a path from v_1 to v_2 in T , $v_1 \cap v_2$ is contained in every node on the path.



(a)



(b)

(a) is a join tree
(for the graph of Fig. 4.14)

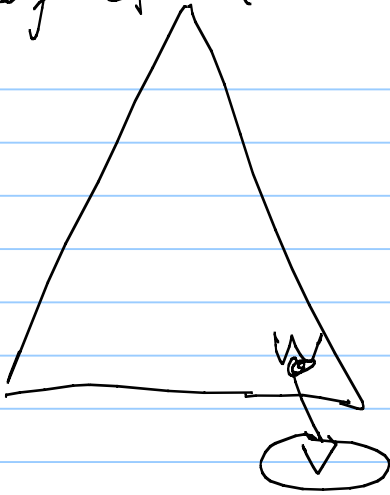
(b) is not a join tree, b/c
 $\{B, C, D, E\} \cap \{B, C, D, G\} = \{B, C, D\}$
 but node DEFI does not
 contain $\{B, C, D\}$!

Theorem 4.3 If the cliques of an undirected graph G can be organized into a join tree, then G is triangulated

Theorem 4.4 If the undirected graph G is triangulated, then G can be organized into a join tree

Proof of Thm 4.3

$$V \setminus W \neq \emptyset$$



Remove every variable in $V \setminus W$. Each of these
can be eliminated w/ out creating fill-ins,
b/c all of their neighbors are in the
closure w . Repeat!

