

- More asymptotic notation
 - Asymptotics of logs, exponentials, $n!$
 - Integral approximations

$f(n) = o(g(n))$ ("little oh")
 means $\forall c > 0, \exists N, \forall n \geq N$
 $f(n) \leq c g(n)$
 (equiv: $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$)
 "f grows strictly slower than g (asymptotically)"

$n = o(n^2)$
 $n \leq c \cdot n^2$ for all $n \geq \frac{1}{c}$

Prf: $n \geq \frac{1}{c} \Rightarrow cn \geq 1$
 $\Rightarrow cn^2 \geq n$

For all $a < b$, ($a, b \in \mathbb{R}$)
 $n^a = o(n^b)$

Fact: $f(n) = o(g(n)) \Rightarrow f(n) = O(g(n))$
 (strict version) (non-strict asymptotic inequality)

$f(n) = O(g(n))$ but
 $f(n) \neq o(g(n))$

Def: $f(n) = \omega(g(n))$ means
 $\forall c > 0, \exists N, \forall n \geq N$
 $f(n) \geq c g(n)$
 (equiv: $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$)
 (equiv: $g(n) = o(f(n))$)

Fix constant $b > 1$. ($b \in \mathbb{R}$)
 $f(n) = b^n$ exponential function to base b
 strictly monotone increasing over $\mathbb{R}$

Fact: for every constant k,
 $n^k = o(b^n)$

Proof: $b = e$, $k \in \mathbb{N}$
 $\lim_{n \rightarrow \infty} \frac{n^k}{e^n} = \lim_{n \rightarrow \infty} \frac{k n^{k-1}}{e^n}$
 $= \dots = \frac{k(n-1) \dots (1)}{e^n} = \lim_{n \rightarrow \infty} \frac{k!}{e^n}$
 $= k! \lim_{n \rightarrow \infty} \frac{1}{e^n} = 0$

$k \in \mathbb{N}: n^k = O(n^{k+1}) = o(e^n)$

$n^k = o(b^n)$
 $n^k \leq c b^n$

$\log_b$ of both sides ($\log_b$ is monotone increasing because $b > 1$)

$k \log_b n \leq \log_b c + n \log_b b$
 $= \log_b c + n$
 $\therefore \log_b n = o(n)$

Fact: For any real $k > 0$,
 $\log_b n = o(n^k)$
 $\lg n = o(n^k)$

Equiv: For any real $l > 0$
 $(\lg n)^l = o(n)$
 (take both sides to the power of $l < \frac{1}{k}$)
 polylog growth

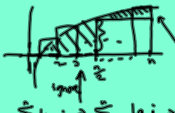
Def: $f(n) = \text{Poly}(g(n))$
 means that $\exists k \geq 0$
 $f(n) = O(g(n)^k)$
 $[n^k = \text{Poly}(n)]$
 "f(n) is polynomially bounded in g(n)"
 $e^n \neq \text{Poly}(n)$

$\Omega(f(n)) = O(g(n))$
 $\circ(g(n))$
 $\Omega(g(n))$
 $\omega(g(n))$

and $h(n)$ is asymptotically positive
 then
 $\Omega(h(n)) = O(g(n)h(n))$
 $\circ(g(n)h(n))$
 $\Omega(g(n)h(n))$
 $\omega(g(n)h(n))$
 respectively.

$\lg n = o(n^\epsilon) \quad \forall \epsilon > 0$
 $n^\epsilon \lg n = o(n^{2+\epsilon}) \quad \forall \epsilon > 0$

$e^x \geq 1+x \quad (\forall x \in \mathbb{R})$
 $e^x \leq 1+x+x^2 \quad (\forall x, |x| \leq 1)$
 $n! = 1 \cdot 2 \cdot 3 \cdots (n-1) \cdot n$
 $= \prod_{i=1}^n i \quad [0! = \prod_{i=1}^0 i = 1]$
 $\lg(n!) = \lg\left(\prod_{i=1}^n i\right) = \sum_{i=1}^n \lg i$
 $\leq \sum_{i=1}^n \lg n = n \lg n$

$\sum_{i=1}^n \lg i \geq \sum_{i=1}^n \lg \frac{i}{2} \geq \sum_{i=1}^n \lg \frac{i}{2}$
 $\geq \frac{n}{2} \cdot \lg\left(\frac{n}{2}\right) = \frac{n}{2} (\lg n - 1)$
 $\therefore \lg(n!) = \Theta(n \lg n)$
 $C_1 n \lg n \leq \lg(n!) \leq C_2 n \lg n$
 for some $0 < C_1 \leq C_2$
 & all sufficiently large n
 Take 2 to the power of both sides:
 $2^{C_1 n \lg n} \leq n! \leq 2^{C_2 n \lg n}$
 $\parallel$
 $(2^{C_1})^{n \lg n} \leq n! \leq (2^{C_2})^{n \lg n}$
 $1 < b < c \Rightarrow b^n = o(c^n)$

Stirling's approximation:
 $n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$
 $\lim_{n \rightarrow \infty} \frac{n!}{\sqrt{2\pi n} \left(\frac{n}{e}\right)^n} = 1$
 $n! = \sqrt{2\pi n} \left(\frac{n}{e}\right)^n e^{\alpha_n}$
 where $\frac{1}{12n+1} < \alpha_n < \frac{1}{12n}$

Integral Approximation
 Assume $a, b \in \mathbb{Z}$ $a \leq b$.

- If $f(x)$ is monotone ascending on $[a-1, b+1]$, then

$$\int_{a-1}^b f(x) dx \leq \sum_{i=a}^b f(i) \leq \int_a^{b+1} f(x) dx$$

- If $f(x)$ is monotone descending on $[a-1, b+1]$, then

$$\int_a^{b+1} f(x) dx \leq \sum_{i=a}^b f(i) \leq \int_{a-1}^b f(x) dx$$

[f is assumed to be Riemann integrable, always true for us]

$\int_1^{n+1} \frac{dx}{x} \leq \sum_{i=1}^n \frac{1}{i} \leq \int_0^n \frac{dx}{x} = \infty$
 $\boxed{H(n+1)} \quad \parallel \quad 1 + \sum_{i=2}^n \frac{1}{i} \leq 1 + \int_1^n \frac{dx}{x} = 1 + \ln n$

$$1 + \int_2^{n+1} \frac{dx}{x} \leq 1 + \sum_{i=2}^n \frac{1}{i}$$

$$\begin{aligned} \mapsto \ln x \Big|_2^{n+1} &= 1 + \ln(n+1) - \ln 2 \\ &= \ln(n+1) + \underbrace{(1 - \ln 2)}_{> 0} \end{aligned}$$

How tight? (monotone ascending case)

U.B - L.B = width of the interval containing

$$\sum_{i=a}^b f(i)$$

||

$$\begin{aligned} \int_a^{b+1} f(x) dx - \int_{a-1}^b f(x) dx \\ &= \int_b^{b+1} f(x) dx - \int_{a-1}^a f(x) dx \\ &\leq \int_b^{b+1} \underline{f(b+1)} dx - \int_{a-1}^a \underline{f(a-1)} dx \\ &= f(b+1) - f(a-1) \end{aligned}$$

Iterated lg function

For $n \geq 1$, define

$$\lg^* n = \text{least } i \text{ such that } \lg^{(i)} n \leq 0$$

!!

$$\underbrace{\lg \cdots \lg \lg \lg n}_{i \text{ times}}$$

$$\lg^* 1 = 1$$

$$\lg^* 2 = 2$$

$$\lg 2 = 1$$

$$\lg \lg 2 = 0$$

$$\lg^* 3 = 3$$

$$\lg^* 4 = 3$$

$$\lg^* 2^{\overbrace{2}^i} = i+1$$

$$\lg(2^{\overbrace{2}^i}) = 2^{\overbrace{2}^{i-1}} \underbrace{\lg 2}_1 = 2^{\overbrace{2}^{i-1}}$$