

Greedy Algos

- Optimization problems
- Optimal substructure
- Recursive solution
- Dynamic programming sol'n?
- Greedy solution

Greedy algo:

Takes a problem and breaks it down into sub or more subproblems that look the best

Set of activities $\{a_1, \dots, a_n\}$

each a_k has a

start time s_k

finish time f_k

duration of a_k is the interval $[s_k, f_k)$

Find largest subset of activities that are compatible (durations don't overlap).

[Sister] [Sister]

- throw out any activity incompatible with a_k

- two sets of remaining activities:

Sister = $\{a_j : f_j \leq s_k\}$

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only subproblems to solve are:

$S_{ij} = \{a_k : f_i \leq s_k \leq f_k \leq s_j\}$

all activities between

a_i and a_j

Add two dummy activities:

$a_0 = [-\infty, 0]$

$a_{n+1} = [0, \infty]$

Solve for S_{ij} $\{0 \leq i, j \leq n+1\}$

Greedy algorithm:

Sort all activities by finish time. Re-index so that $f_1 \leq f_2 \leq f_3 \leq \dots \leq f_n$

While there are activities left

- choose activity with earliest

start finish time

- transfer it to the solution

- remove all activities incompatible with this choice.

$a_1 \rightarrow$ solution
remove all a_i st. $s_i < f_1$
(incompatible with a_1)

Don't need a table of sol'n's to sub problems. Usually a single recursive algo is fine.

Huffman coding

Text file of chars from some alphabet $C = \{a, b, c, d, \dots\}$

Encode in binary as efficiently (space) as possible

(optimal compression)

assuming that each symbol

is encoded by a binary string.

Prefix code:

$s_i \rightarrow b_i$ $b_1, \dots, b_n \in \{0, 1\}^*$

$s_i \rightarrow b_i$ b_i code words

$s_i \rightarrow b_i$

such that no b_i is a prefix of b_j for any $i \neq j$.

$a \rightarrow 0$

$b \rightarrow 10$

$c \rightarrow 11$

Binary encoding tree

(gives a prefix code)

$C = \{a, b, c, d\}$

File F has 1000 a's

250 b's

500 c's

250 d's

$a \rightarrow 00$

$b \rightarrow 01$

$c \rightarrow 10$

$d \rightarrow 11$

Encoded file has 4000 bits

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1000 a's $a \rightarrow 0$
 250 b's $b \rightarrow 110$
 500 c's $c \rightarrow 10$
 250 d's $d \rightarrow 111$



encoded file has
 $\frac{1000 \times 1}{a's} + \frac{250 \times 3}{b's} + \frac{500 \times 2}{c's} + \frac{250 \times 3}{d's}$

$$\begin{array}{r} 1000 \\ 750 \\ 1000 \\ 750 \\ \hline 3500 \end{array} < 4000$$

Opt problem:
 given an alphabet

$$C = \{s_1, \dots, s_n\}$$

each letter s_i has frequency f_i

find an optimal ^{binary} prefix code.
 (Huffman code)

given a code

$s_i \rightarrow b_i$ let l_i be
 $\vdots$ length of b_i
 $s_k \rightarrow b_k$

$$\text{cost} = \sum_{i=1}^n l_i f_i = \text{size of encoded file}$$

we want to choose a prefix code to minimize this.

Greedy solution: builds an optimal encoding tree.

Huff(C) $x \in C$

$x.\text{freq} \leftarrow \text{freq of } x$

// C is nonempty

if $|C| = 1$

return a leaf with the one char of C as root

else

let $x, y \in C$ be two letters with minimum frequency

introduce a new char

$$z = \begin{array}{c} x \quad y \\ \diagup \quad \diagdown \\ \end{array} \text{ not in } C$$

$$z.\text{freq} := x.\text{freq} + y.\text{freq}$$

remove x, y from C,

add z to C, forming

a new alphabet C'

return Huff(C').

Nonrecursive solution

while C has ≥ 2 letters

- remove letters x, y from C with min frequency

- form new letter $z = \begin{array}{c} x \quad y \\ \diagup \quad \diagdown \\ \end{array}$

with $z.\text{freq} = x.\text{freq} + y.\text{freq}$

- add z to C

end while

return sole char of C (tree)



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