

Quick solution
Proof sketch of Master Thm

$$T(n) = \Theta(n^2)$$

Master Method:

$$T(n) = 2T\left(\frac{n}{10}\right) + n^2$$

$$a = 2$$

$$b = 10, \quad t = \log_b a$$

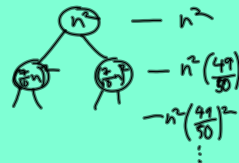
$$f(n) = n^2 = \log_b 2 \cdot \sqrt{2}$$

Case 3:

$$T(n) = \Theta(f(n)) \because \left(\frac{10}{9}\right)^2 = 2$$

$$= \Theta(n^2)$$

Tree method:



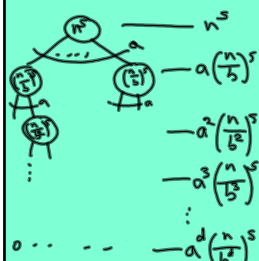
$$\begin{aligned} n^2 \leq T(n) &\leq \sum_{i=0}^{\log_{10} n} n^2 \left(\frac{1}{10}\right)^i \\ &= n^2 \sum_{i=0}^{\log_{10} n} \left(\frac{1}{10}\right)^i \\ &\leq n^2 \sum_{i=0}^{\infty} \left(\frac{1}{10}\right)^i \\ &= n^2 \left(\frac{1}{1 - \frac{1}{10}}\right) \\ &= n^2(50) \end{aligned}$$

$$T(n) = \Theta(n^2)$$

Sketch proof of Master Thm using tree method

Special case: $f(n) = n^s$
(constant $s \geq 0$).

$$T(n) = aT\left(\frac{n}{b}\right) + n^s$$



where $d = \text{depth of the tree}$
 $d = \log_b n$

$$\begin{aligned} T(n) &= \sum_{i=0}^d a^i \left(\frac{n}{b^i}\right)^s \\ &= n^s \sum_{i=0}^d \frac{a^i}{b^{is}} \\ &= n^s \sum_{i=0}^d \left(\frac{a}{b^s}\right)^i \\ &= n^s \left(\sum_{i=0}^d r^i\right) \text{ where } r = \frac{a}{b^s} \end{aligned}$$

Case 2: $r = 1$:

$$b^s = a \quad \log_b \text{ of both sides}$$

$$s = \log_b a = \log_b a = t$$

(Case 2 of Master Thm)

$$T(n) = n^s(d+1)$$

$$= \Theta(n^s d)$$

$$= \Theta(n^s \log_b n)$$

$$= \Theta(n^s \lg n)$$

all logs equal.

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$$T(n) = \sum_{i=0}^d r^i \quad r = \frac{a}{b}$$

Case 3: $r < 1$

$$\frac{a}{b^s} < 1 \rightarrow a < b^s$$

$$t = \log_b a < s$$

(case 3 of Master-Thm)

$$n^s \leq T(n) = \sum_{i=0}^d r^i$$

$$\leq \sum_{i=0}^{\infty} r^i = O(n^s)$$

$\frac{1}{1-r} > 0$

$$\therefore T(n) = \Theta(n^s) = \Theta(f(n))$$

tree is top-heavy.

Case 1: $r > 1$.

$$T(n) = \sum_{i=0}^d r^i$$

$$= n^s \left(\frac{r^{d+1} - 1}{r - 1} \right)$$

$$= \Theta(n^s r^{d+1})$$

$$= \Theta(n^s r^d)$$

$$d = \log_b n$$

$$= \Theta(n^s r^{\log_b n})$$

$$= \Theta(n^s n^{\log_b r})$$

$$= \Theta(n^{s + \log_b r})$$

$$\left[r = \frac{a}{b^s} \right]$$

$$= \Theta(n^{s + \log_b \left(\frac{a}{b^s} \right)})$$

$$= \Theta(n^{s + \log_b a - \frac{\log_b b^s}{s}})$$

$$= \Theta(n^{\log_b a})$$

$$= \Theta(n^t)$$

tree is bottom-heavy.

$$- n^s$$

$$- r n^s$$

$$- r^2 n^s$$

Counting things
(Appendix C)

A, B, disjoint

|A| ways to select from A

+

|B| B

|A| + |B| possible ways to select an element from either A or B
(sum rule)

(product rule) A, B arbitrary

|A| · |B| ways to select a pair (a, b) where a ∈ A and b ∈ B

$$|A \times B|$$

labeled k elements from A (1, 2, 3, ..., k)

with replacement

|A|^k ways (product rule)

A, B sets

$$|A| = n$$

$$|B| = m$$

How many functions
from A into B?

m^n many ways

$$A = \{a_1, a_2, \dots, a_n\}$$

functions = # n-tuples from B
product rule says $= |B|^n = m^n$

"without replacement"

A select a sequence of k elements
from A, all distinct.

$$n(n-1)(n-2) \dots (n-k+1)$$

k factors

$$= \frac{n!}{(n-k)!}$$

$$k = n: \frac{n!}{(n-n)!} = \frac{n!}{0!} = n!$$

$n! =$ # permutations of A
 $=$ # possible ways of ordering
the elements of A

A, B $|A| = n, |B| = m$

How many 1-1 functions
from A into B?

$$= \frac{n!}{(n-m)!} \quad (m \geq n)$$

(if $n < m$, then $\emptyset$)

A permutation of A is
a 1-1 function of A onto A
($n!$ of these)

$$|A| = n$$

k elements of A, without
regard to order.

counting the # subsets of A
with k elements.

$$\frac{n!}{(n-k)! k!} = \binom{n}{k} \quad \text{"n choose k"}$$

If order does matter, then
we saw $\frac{n!}{(n-k)!}$ many ways;

this counts the same
collection of k elements
exactly $k!$ many times
(one for each ordering of the
 k elements).
So,

$$\frac{n!}{(n-k)! k!} = k! \left(\# \text{ of } k\text{-subsets of A} \right)$$

$$\left(\# \text{ of } k\text{-subsets of } n \text{ elements} \right) = \frac{n!}{k! (n-k)!} = \binom{n}{k}$$

$$0 \leq k \leq n$$

Binomial Thm:

For any numbers x, y and
integer $n \geq 0$

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k}$$

So in particular

$$x=y=1:$$

$$\sum_{k=0}^n \binom{n}{k} = 2^n$$

$$x=1, y=-1$$

$$\binom{n}{0} - \binom{n}{1} + \binom{n}{2} - \binom{n}{3} + \dots = 0$$

(if $n > 0$).