

This document contains slides from the lecture, formatted to be suitable for printing or individual reading, and with some supplemental explanations added. It is intended as a supplement to, rather than a replacement for, the lectures themselves — you should not expect the notes to be self-contained or complete on their own.

1 Introduction

CLRS 14

In many cases, we can use existing data structures in unexpected ways by **augmenting** the data they store.

Basic steps:

1. **Choose** an underlying data structure.
2. **Determine** additional information to maintain in that data structure.
3. **Modify** the operations of that data structure to maintain that information.
4. **Develop** new operations using that information.

This approach is much more common than “starting from scratch” with a new data structure.

Many examples are based on balanced search trees.

2 Dynamic order statistics

Suppose we want a data structure that supports these operations:

- INSERT(k)
- SEARCH(k)
- DELETE(k)
- SELECT(i) — find the i^{th} smallest element
- RANK(v) — how many elements are smaller than the one at node v ?

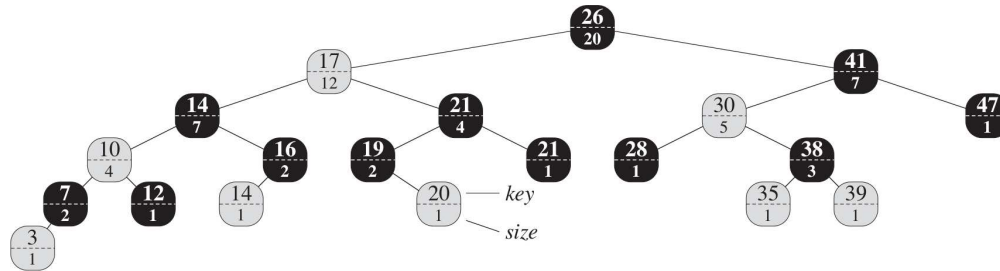
This differs from the standard selection problem, because the set is **dynamic** — elements may be added or deleted.

We can form a data structure that supports these operations by augmenting your favorite rotation-based balanced binary search tree data structure.

3 Order statistic trees

In addition to the usual attributes (key, left, right, parent), add a new attribute:

- Store the number of nodes in the subtree rooted at v as $v.size$.

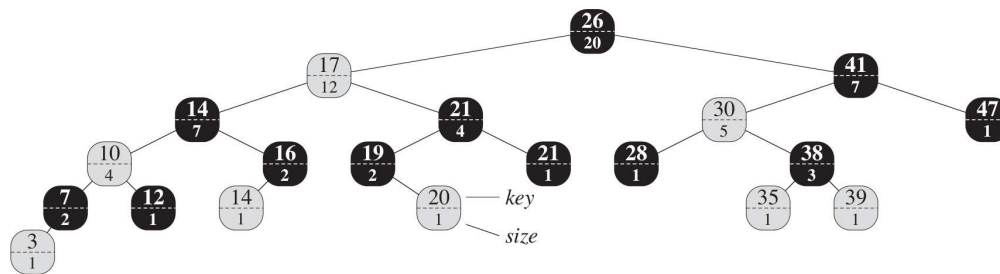


(CLRS Fig. 14.1)

$$v.size = v.left.size + v.right.size + 1$$

4 Maintaining the size attribute: INSERT

For each INSERT: Increment the size of each node along the way.



5 Maintaining the size attribute: ROTATE

Each rotation changes the size for only two nodes: The two nodes incident to the edge being rotated.

6 SELECT in Order Statistic Trees

The SELECT operation in an order statistic tree looks much like QUICKSELECT:

```
OST-SELECT( $v, i$ )
 $r = v.left.size + 1$ 
if  $i = r$  then
    return  $v$ 
else if  $i < r$  then
    return OST-SELECT( $v.left, i$ )
else
    return OST-SELECT( $v.right, i - r$ )
end if
```

This takes time proportional to the height of the tree, which is $\Theta(\lg n)$ for a balanced binary search tree.

7 *Computing rank in order statistic trees*

How can we use this data structure to compute the **rank** of a given node in the tree?

(CLRS 342)