- An interpretation is an assignment of values to all variables.
- A model is an interpretation that satisfies the constraints.
- Often we don't want to just find a model, but want to know what is true in all models.
- A proposition is statement that is true or false in each interpretation.

- Specifying logical formulae is often more natural than filling in tables
- It is easier to check correctness and debug formulae than tables
- We can exploit the Boolean nature for efficient reasoning
- We need a language for asking queries (of what follows in all models) that may be more complicated than asking for the value of a variable
- It is easy to incrementally add formulae
- It can be extended to infinitely many variables with infinite domains (using logical quantification)

Step 1 Begin with a task domain.

Step 2 Choose atoms in the computer to denote propositions. These atoms have meaning to the KB designer.

Step 3 Tell the system knowledge about the domain.

Step 4 Ask the system questions.

— the system can tell you whether the question is a logical consequence.

 You can interpret the answer with the meaning associated with the atoms.

In computer:

In user's mind:

- *light1_broken*: light #1 is broken
- *sw_up*: switch is up
- *power*: there is power in the building
- unlit_light1: light #1 isn't lit
- *lit_light*2: light #2 is lit

Conclusion: *light1_broken*

- The computer doesn't know the meaning of the symbols
- The user can interpret the symbol using their meaning

- An atom is a symbol starting with a lower case letter
- A body is an atom or is of the form $b_1 \wedge b_2$ where b_1 and b_2 are bodies.
- A definite clause is an atom or is a rule of the form $h \leftarrow b$ where h is an atom and b is a body.
- A knowledge base is a set of definite clauses

- An interpretation / assigns a truth value to each atom.
- A body $b_1 \wedge b_2$ is true in I if b_1 is true in I and b_2 is true in I.
- A rule h ← b is false in I if b is true in I and h is false in I. The rule is true otherwise.
- A knowledge base *KB* is true in *I* if and only if every clause in *KB* is true in *I*.

- A model of a set of clauses is an interpretation in which all the clauses are *true*.
- If KB is a set of clauses and g is a conjunction of atoms, g is a logical consequence of KB, written $KB \models g$, if g is true in every model of KB.
- That is, $KB \models g$ if there is no interpretation in which KB is *true* and g is *false*.

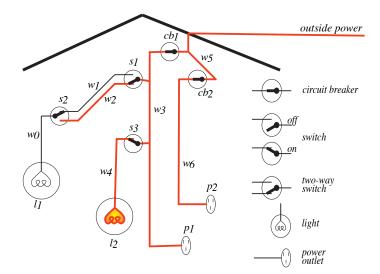
$$\mathcal{KB} = \left\{ egin{array}{c} p \leftarrow q. \\ q. \\ r \leftarrow s. \end{array}
ight.$$

	p	q	r	5			
I_1	true		true	true	is a model of <i>KB</i>		
I_2	false	false	false	false	not a model of <i>KB</i>		
I_3	true	true	false	false	is a model of <i>KB</i>		
<i>I</i> 4	true	true	true	false	is a model of <i>KB</i>		
<i>I</i> 5	true	true	false	true	not a model of <i>KB</i>		
$\mathit{KB}\models \mathit{p}, \mathit{KB}\models \mathit{q}, \mathit{KB} eq r, \mathit{KB} eq s$							

- 1. Choose a task domain: intended interpretation.
- 2. Associate an atom with each proposition you want to represent.
- 3. Tell the system clauses that are true in the intended interpretation: axiomatizing the domain.
- 4. Ask questions about the intended interpretation.
- 5. If $KB \models g$, then g must be true in the intended interpretation.
- 6. Users can interpret the answer using their intended interpretation of the symbols.

- The computer doesn't have access to the intended interpretation.
- All it knows is the knowledge base.
- The computer can determine if a formula is a logical consequence of KB.
- If $KB \models g$ then g must be true in the intended interpretation.
- If $KB \not\models g$ then there is a model of KB in which g is false. This could be the intended interpretation.

Electrical Environment



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Representing the Electrical Environment

$light_1$.	$\textit{lit_l_1} \gets \textit{live_w_0} \land \textit{ok_l_1}$		
light_b.	$\mathit{live_w_0} \leftarrow \mathit{live_w_1} \land \mathit{up_s_2}.$		
$down_{s_1}$.	$\textit{live_w_0} \leftarrow \textit{live_w_2} \land \textit{down_s_2}.$ $\textit{live_w_1} \leftarrow \textit{live_w_3} \land \textit{up_s_1}.$		
up_s ₂ .			
up_5 ₃ .	$\mathit{live}_w_2 \leftarrow \mathit{live}_w_3 \land \mathit{down}_s_1.$		
ok_l ₁ .	$lit_{-}l_{2} \leftarrow live_{-}w_{4} \land ok_{-}l_{2}.$		
ok_h.	$live_w_4 \leftarrow live_w_3 \wedge up_s_3.$		
ok_cb1.	$live_p_1 \leftarrow live_w_3.$		
ok_cb ₂ .	$live_w_3 \leftarrow live_w_5 \land ok_cb_1.$		
live outside.	$live_p_2 \leftarrow live_w_6.$		
inve_outside.	<i>live_w</i> ₆ \leftarrow <i>live_w</i> ₅ \land <i>ok_cb</i> ₂ .		
	$live_w_5 \leftarrow live_outside.$		