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Note Title

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Normal Forms

Claim. If x is not free in G , then

$$Q_{[x/u]}(G) = Q(G)$$

Proof. $Q_{[x/u]}$ is equal to Q except

that x is assigned u .

If x is not free, then, regardless of the assignment of x (either Q or $Q_{[x/u]}$), x is assigned

according to rule for \forall or the rule for \exists .

So, the change in mapping for x does not matter. \square

$$\text{Exercise 55} \quad (\forall x F \vee \forall x G) \neq \forall x (F \vee G) \quad H(x) = F$$

$$(\exists x F \wedge \exists x G) \neq \exists x (F \wedge G) \quad J(x) = G$$

The suitable structure (interpretation) provides the required counterexample:

$$U_a = \{a, b\} \quad \mathcal{A}(H) (= H^a) = \{a\} \quad \mathcal{A}(J) (= J^a) = \{b\}$$

$$Q(\forall x H(x) \vee \forall x J(x)) = 0, \text{ but}$$

$$Q(\forall x (H(x) \vee J(x))) = 1, \text{ and}$$

$$Q(\exists x H(x) \wedge \exists x J(x)) = 1, \text{ but}$$

$$Q(\exists x (H(x) \wedge J(x))) = 0$$

Exercise 56 Show that $F = (\exists x P(x) \rightarrow P(y))$
is equivalent to $G = \forall x (P(x) \rightarrow P(y))$.

[$P(x) \leftarrow Q(x, y)$ in A1 Log]

$\forall x \forall y P(x) \leftarrow Q(x, y)$

$\forall x (P(x) \stackrel{\equiv}{\leftarrow} \exists y Q(x, y))$]

Let $\mathcal{Q} = (U_{\mathcal{Q}}, \mathcal{I}_{\mathcal{Q}})$ a suitable structure (interpretation) for F and G .

$$\mathcal{Q}(\exists x P(x) \rightarrow P(y)) = 1$$

iff $\mathcal{Q}(\sim \exists x P(x) \vee P(y)) = 1$ (defn. of implication)

iff either $\mathcal{Q}(\sim \exists x P(x)) = 1$ or $\mathcal{Q}(P(y)) = 1$ (4. on p. 47)

iff $\mathcal{Q}(\forall x \sim P(x))$ or $\mathcal{Q}(P(y)) = 1$ (1b p. 52)

iff $\mathcal{Q}(\forall x \sim P(x) \vee P(y)) = 1$ (4 on p. 47)

iff $\mathcal{Q}(\forall x (\sim P(x) \vee P(y))) = 1$ (2b p. 52 with $F = \sim P(x)$
and $G = P(y)$)

iff $\mathcal{Q}(\forall x (P(x) \supset P(y))) = 1$ (defn of implication);

It would be shorter to provide \mathcal{Q} directly.

